

2.4.2.

$$\phi \stackrel{\text{def}}{=} \forall x \exists y \exists z \left(P(x, y) \wedge P(z, y) \wedge \left(P(x, z) \rightarrow P(z, x) \right) \right)$$

(a) $A \stackrel{\text{def}}{=} \mathbb{N}$

$$P \stackrel{\text{def}}{=} \{ (m, n) \mid m < n \}$$

$$m P n \stackrel{\text{def}}{=} m < n$$

Abusing notation:

$$\forall x \in \mathbb{N} \exists y \in \mathbb{N} \exists z \in \mathbb{N}$$

$$\left(x < y \wedge z < y \wedge \left(x < z \rightarrow z < x \right) \right)$$

$$M \models \phi$$

2.4.3.

$$\forall x \neg P(x, x)$$

$$A = \{ \text{apple} \}$$

$$P^{\mathcal{M}_1} = \{ \}$$

$$\mathcal{M}_1 \models \forall x \neg P(x, x)$$

$$P^{\mathcal{M}_2} = \{ (\text{apple}, \text{apple}) \}$$

$$\mathcal{M}_2 \not\models \forall x \neg P(x, x)$$

2.4.12.

$$(a) \quad \forall x \forall y (S(x,y) \rightarrow S(y,x)) \rightarrow \forall x \neg S(x,x) = \phi$$

$$A = \{ \text{apple} \}$$

$$S^{M_2} = \{ (\text{apple}, \text{apple}) \}$$

$$M_2 \neq \phi$$

$$A = \{ \text{apple} \}$$

$$S^{M_1} = \{ \}$$

$$M_1 \models \phi$$

$$A = \mathbb{N}$$

$$n \in S^{M_3}_m \Leftrightarrow n = m$$

$$M_3 \neq \phi$$